Design considerations for math libraries

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Context

- My primary experience with FFTs is via MiniDFT benchmarking 2016-2017.
- I've worked on NWChem continuously since 2006, with a focus on tensor contractions (i.e. not FFTs).
- The higher-level, distributed structure of all math libraries are roughly the same.

Bad Things

- Synchronization
- Unnecessary data movement
- Synchronization
- Load-imbalance
- Synchronization
- Assuming all 50+ cores should be used for computation.
- Synchronization

GPUs improve the situation quite a bit because asynchronous execution is normative, unlike CPU execution.

Bottom-up tensor contractions

Z(d,b,a,c) = X(b,f,a,e) * Y(d,e,c,f) // Einstein notation

X2(a,b,e,f) = X(b,f,a,e) // Transpose O(N⁴)

Y2(e,f,c,d) = Y(d,e,c,f) // Transpose O(N⁴)

Z2(a,b,c,d) = X2(a,b,e,f) * Y2(e,f,c,d)

Z(d,b,a,c) = Z2(a,b,c,d)

// Matrix multiplication O(N⁶)

// Transpose O(N⁴)

https://github.com/jeffhammond/spaghetty/blob/main/papers/old/spaghetty_chapter_v1.pdf

Bottom-up distributed tensor contractions

```
FORALL a,b,c,d,e,f:
Fetch X(b,f,a,e)
X2(a,b,e,f) = X(b,f,a,e)
Fetch Y(d,e,c,f)
Y2(e,f,c,d) = Y(d,e,c,f)
Z2(a,b,c,d) = X2(a,b,e,f) * Y2(e,f,c,d)
Z(d,b,a,c) = Z2(a,b,c,d)
Update Z(d,b,a,c)
```

// index sets group in tiles // Communication $O(N^4)$ // Transpose O(N⁴) // Communication $O(N^4)$ // Transpose O(N⁴) // Matrix multiplication $O(N^6)$ // Transpose O(N⁴) // Communication O(N⁴)

FORALL a,b,c,d,e,f: // index sets group in tiles // Communication $O(N^4)$ NB Fetch X(b,f,a,e) // Communication $O(N^4)$ NB Fetch Y(d,e,c,f) Wait on X. Y X2(a,b,e,f) = X(b,f,a,e)// Transpose O(N⁴) Y2(e,f,c,d) = Y(d,e,c,f)// Transpose O(N⁴) // Matrix multiplication $O(N^6)$ Z2(a,b,c,d) = X2(a,b,e,f) * Y2(e,f,c,d)Wait for 7 to be available // Transpose O(N⁴) Z(d,b,a,c) = Z2(a,b,c,d)// Communication $O(N^4)$ NB Update Z(d,b,a,c)

Composing math and communication

- Everything should be asynchronous but not tedious to manage
 - Explicit events/requests are tedious.
 - Completion callbacks are okay, but how far can we recurse? And what's the MPI situation?
 - CUDA streams work but streams are an implementation tool not the right semantic.
 - CUDA graphs and other graph-based execution models make sense.
 - Explicit dependencies ala OpenMP 4+ are pretty good, but lack an object model.

- There are known gaps with MPI here.
 - Completion callbacks are not supported (as in DCMF/PAMI).
 - Active messages are not supported directly (not that they are useful in FFTs).
 - Driving MPI with a progress thread that manages completions is often reasonable.

Top-down tensor contractions

Z(d,b,a,c) = X(b,f,a,e) * Y(d,e,c,f)

Z["dbac"] = X["bfae"]*Y["decf"];

// Einstein notation

// actual C++ code for CTF

- 1) This is the only known way to build formally optimally implementations.
- 2) Everyone has to agree on the need for this solution and the notation.
- 3) This is fully synchronous at the term level and higher-level task parallelism is not supported by this interface.

Suggestions for FFT people

- Higher-level interfaces will do a much better job adapting to novel memory hierarchies.
- Come up with the high level library interface that makes everyone happy, but make it asynchronous like a nonblocking collective.
 - Apparently, this is hard because use cases pretty wildly across quantum and non-quantum use cases.
 - Nonblocking != asynchronous. Must provide resources to drive background progress.
- Build FFT kernel components that fit into asynchronous tasking systems that allow composition of computation, communication, allocation, etc.
 - No asynchronous tasking system for Fortran except OpenMP 4.5 (tasking is CPU-only).